

The Polyhedral Compilation Package (PCP)

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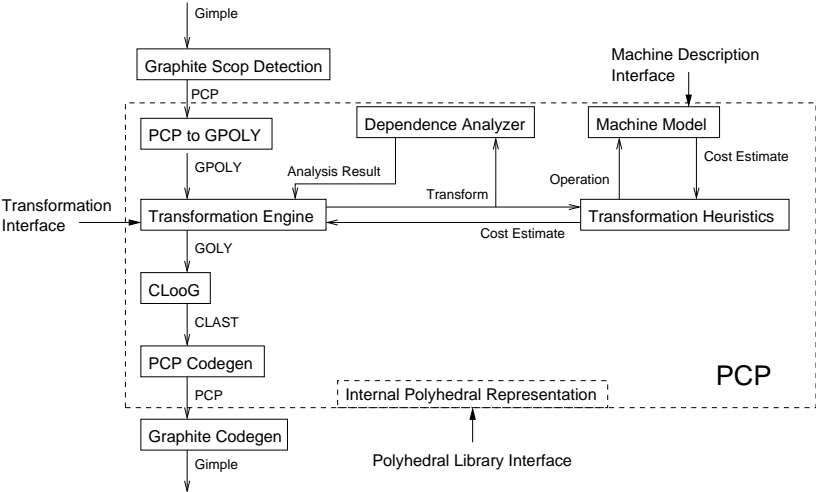
Outline

- ▶ Motivation
- ▶ PCP framework
- ▶ PCP language
- ▶ Annotations
- ▶ Test infrastructure
- ▶ Loop interface
- ▶ Machine model and heuristics

Motivation for PCP

- ▶ Graphite was created to improve loop optimizations in GCC
- ▶ PCP defines a set of interfaces to make the code more modular by separating GIMPLE and the polyhedral representation (GPOLY)
- ▶ Well defined interfaces that are easy to use and understand
- ▶ Ability to evolve the framework independently of GIMPLE
- ▶ Test the framework without the need for a front end
- ▶ Source-to-source translation is desirable to allow iterative optimization

Graphite/PCP Overview



Design Criteria for the PCP Language

- ▶ Easy to understand
 - ▶ Imperative language constructs that easily map from and to GIMPLE.
- ▶ Small
 - ▶ Design a language that is only as expressive as the polyhedral model.
 - ▶ Cannot represent all possible operators/statements that exist: use black boxes and only expose array reads/writes. *Mapping IR to a black box is not trivial in some cases.*
- ▶ Easy to integrate and extend
 - ▶ Allow auxiliary information to be passed to and from PCP.
- ▶ Testable and Debuggable
 - ▶ Test framework
 - ▶ Define syntax for the PCP language

PCP Language Overview

- ▶ Imperative language that expresses data communication via array accesses.
- ▶ Compilation unit is a program fragment of structured code that can be expressed with simple loops, conditional statements and array accesses. This is called a SCoP (Static Control Part).
- ▶ Computations are expressed as black boxes, only the array accesses are represented.
- ▶ Expressions for array accesses and conditions must be linear combinations of induction variables and parameters.

Language Components

- ▶ Type: `array(Dimension0, ..., DimensionN)`
- ▶ Variable: `variable(ArrayType)`
- ▶ Parameter: `parameter()`
- ▶ Linear expressions: `i, N, 42, +(*(4, j), 17)`
- ▶ Boolean expressions: `ge(N, i), and(ge(j,5) ge(i, *(2, j))),`
- ▶ Array access:
`use/def(Array, Subscript0, ..., SubscriptN)`
- ▶ Statements:
 - ▶ `copy(Def, Use)`
 - ▶ `stmt(Access0, ..., AccessN)`
 - ▶ `guard(Condition) { Body }`
 - ▶ `loop(Iv, Start, Condition, Stride) { Body }`
- ▶ SCoP: `scop(Inputs, Outputs, Parameters) { Body }`

Example

C Code (assume a/b/c[1000][1000])

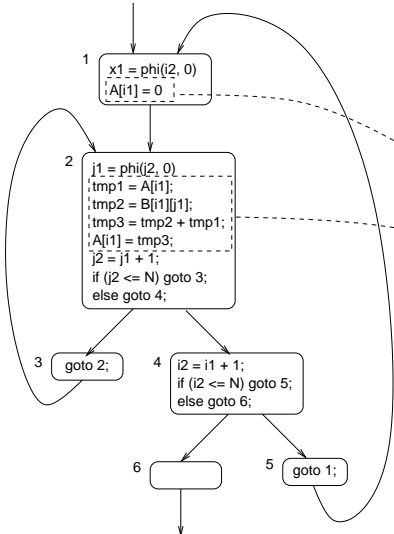
```
for (int i = 0; i < n; i++)
  for (int j = 0; j < 100; j++)
    A[i][j] = B[j][i] + C[i+1][i+j];
```

PCP Code:

```
N <- parameter()
A <- variable(array(1000, 1000))
B <- variable(array(1000, 1000))
C <- variable(array(1000, 1000))
scop (inputs(B,C), outputs(A), parameters(N))
{
  loop(i <- iv(), 0, ge(N,i), 1)
  {
    loop(j <- iv(), 0, ge(100,j), 1)
    {
      // stmt1 maps to the add and assignment
      stmt1(def(A, i, j), use(B, j, i),
            use(C, +(i, 1), +(i, j)))
    }
  }
}
```


Gimple to PCP

GIMPLE



PCP

```
N <- parameter()
A <- variable(array(N))
B <- variable(array(N,N))
scop(inputs(A), outputs(B), parameters(N))
{
  loop(i <- iv(), 0, ge(N, i), 1)
  {
    stmt1(def(A, i))
    loop(j <- iv(), 0, ge(N, j), 1)
    {
      stmt2(def(A, i), use(A, i), use(B, i, j))
    }
  }
}
```



PCP to GPOLY

PCP

```
N ← parameter()
A ← variable(array(N))
B ← variable(array(N,N))
scop(inputs(A), outputs(B), parameters(N))
{
0 i loop(i ← iv(), 0, ge(N, i), 1)
  {
0 stmt1(def(A, i))
1 j loop(j ← iv(), 0, ge(N, j), 1)
  {
0 stmt2(def(A, i), use(A, i), use(B, i, j))
  }
  }
}
```

GPOLY

```
stmt1(def(A, i))
```

Domain	Schedule
(i >= 0)	(0,i,0)
(N >= i)	

```
stmt2(def(A, i), use(A, i), use(B, i, j))
```

Domain	Schedule
(i >= 0)	(0,i,1,j,0)
(N >= i)	
(j >= 0)	
(N >= j)	

Annotations

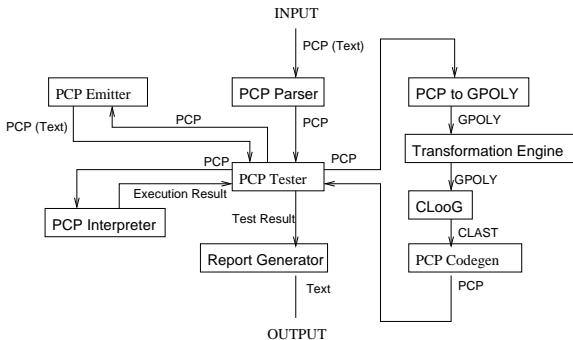
- ▶ Annotate any object in the PCP ASTs to express extra information
- ▶ Does not encode information that is needed to ensure correctness
- ▶ Parser and emitter can handle any annotation since they are purely syntactic

Examples:

```
copy(def(A, i), use(B, i) | lineinfo("test.pcp", 17))  
stmt(def(X, i, j), use(Y, i, j) | costs(size(11),  
                                     time(120)));
```

```
loop(i <- iv(), 0, ge(N, i), 1 | parallel())
```

Test Infrastructure



- ▶ Test cases in the form of PCP code with annotations
- ▶ Easy to dump out a SCoP to a file from any compiler
- ▶ Syntactic tests done by diffing with expected result.
- ▶ Semantic tests are either static tests using analyzes, or dynamic tests that will use an interpreter.
- ▶ Interpreter output will be a trace of memory accesses.

Test Case Example: Loop fusion

```
loop1 <- loop(i <- iv(), 0, ge(N, i), 1))
{
  stmt1(def(B, i, j), use(A, i, j));
}

loop(j <- iv(), 0, ge(N, j), 1 | fusable(loop1))
{
  stmt2(def(C, i, j), use(B, i, j));
}
```

- ▶ If a test fails the file name and line number of the PCP file can be reported.
- ▶ For generic testing the annotations can be ignored and a test can instead be specified via a flag to the tester.
- ▶ May still need a fancier pretty printer that can dump HTML with links/scripts for fast debugging.

Classical Loop Transformations Interface

PCP exposes a classical loop transform interface that can be used to drive the transformations that PCP applies. The interface is based on annotations that are set on the PCP trees:

- ▶ `loop1 (... | fuse (loop2))`
- ▶ `stmt1 (... | move (stmt2))`
- ▶ `loop1 (... | skew (factor))`
- ▶ `loop1 (... | shift (offset))`
- ▶ `loop1 (... | interchange (loop2))`
- ▶ `loop1 (... | stripMine (factor))`
- ▶ `loop1 (... | unroll (factor))`
- ▶ `loop1 (... | reverse)`
- ▶ `loop1 (... | parallelize)`

Machine description and costs for heuristics

Work still in progress but we know we need the following:

- ▶ Description of the machine/system that the code should be optimized for:
 - ▶ Memory hierarchy (caches, latencies, bandwidth)
 - ▶ Number of processors/cores
 - ▶ Vector operation costs
- ▶ Annotations from GCC to estimate code size and execution time.

Conclusion

PCP provides

- ▶ a language for an imperative language restricted to the constraints of the polyhedral model
- ▶ a classical loop transform interface
- ▶ a test infrastructure
- ▶ modularity and ease of debugging

A polyhedral representation GPOLY is contained in PCP